

VITERBI BIT DETECTION METHOD AND DEVICE

BACKGROUND OF THE INVENTION

Field Of The Invention

[0001] The present invention relates to a row-based Viterbi bit detection method for detecting the bit values of bits of a channel data stream stored on a record carrier. Further, the present invention relates to a corresponding row-based Viterbi bit detector, a method of reproduction of a user data stream, a corresponding reproduction device and a computer program for implementing said methods. In particular, the present invention relates to a row-based Viterbi bit detection method for information written in a two-dimensional way on a record carrier, such as an optical disc or a memory card. The present invention could also be regarded as relating to Partial Response Maximum-Likelihood (PRML) bit detection, i.e., the invention also relates to a PRLM bit detection method and device.

Description Of The Related Art

[0002] ~~European patent~~ Patent application ~~Application No.~~ 01203878.2 discloses a method and system for multi-dimensionally coding and/or decoding an information to/from a lattice structure representing channel bit positions of said coded information in at least two dimensions. Encoding and/or decoding is performed by using a quasi close-packed lattice structure. For the case of

three-dimensional encoding and/or decoding, preferably a (quasi) hexagonally close packed (hcp) lattice structure is to be used. Another possibility in three dimensions is the use of a (quasi) face-centered cubic (fcc) lattice structure. For the case of two-
5 dimensional encoding and/or decoding, preferably, a quasi-hexagonal lattice structure is to be used. Another possibility in two dimensions could be the use of a quasi square lattice structure. For the sake of a more simple and clear description of the object of the present invention, special attention is given to the two-
10 dimensional case. The higher-dimensional cases can be derived as more or less straightforward extensions of the two-dimensional case. The special situation of the one-dimensional case which comprises only a single row of bits, boils down to the very classical case of PRML bit detection as is well known in the state
15 of the art for one-dimensional modulation and coding, as, for instance, described in Chapter 7 "Viterbi Detection" by Jan Bergmans, "Digital Baseband Transmission and Recording", Kluwer Academic Publishers, 1996.

[0003] In one-dimensional recording on optical discs, the
20 channel bits of the channel data stream are recorded along a spiral track, the spiral being one bit wide. For two-dimensional recording, the channel bits of a channel data stream can also be recorded along a spiral, albeit a broad spiral, that consists of a number of bit rows which are aligned with respect to each other in
25 the radial direction, that is, in the direction orthogonal to the

spiral direction. The additional alignment of bit rows can also be obtained in another direction not strictly orthogonal to the spiral direction, but in a direction making a certain non-zero angle with the spiral direction.

5 [0004] A PRML bit detection apparatus for deriving a bit sequence from an input information signal is disclosed in WO 00/18016, corresponding to U.S. Patent 6,580,766. The apparatus comprises input means for receiving the input information signal, sampling means for sampling the input information signal at
10 sampling instants so as to obtain samples of the input information signal at said sampling instants, conversion means for converting an array of said samples into an array of bits of a first or a second binary value, detection means for repeatedly detecting a state for subsequent sequences of n subsequent bits of said array
15 of bits, said subsequent sequences being obtained by shifting a time window of n subsequent bits each time over one bit in time, means for establishing the best path through the states, and deriving means for deriving a sequence of bits in accordance with the best path through said states. In that apparatus, n is larger
20 than 3, and sequences of n subsequent bits having $n-1$ directly successive bits of the same binary value are allocated to the same state. In a specific embodiment, n is an odd number larger than 4. In that specific embodiment, sequences of n subsequent bits having $n-2$ directly successive bits of the same binary value as the
25 central $n-2$ bits in such n -bit sequence, are allocated to the same

state. This results in a PRML detection apparatus with reduced complexity.

[0005] A full-fledged PRML bit detector for 2D bit-arrays would require a trellis which is designed for the complete width of the broad spiral, with the drawback of an enormous state-complexity that leads to a completely impractical algorithm, since it cannot by far be implemented even in the fastest hardware of the coming decennia.

[0006] A 2D PRML bit detector is disclosed in "Study of Recording Methods for Advanced Optical Disks", S. Taira, T. Hoshizawa, T. Kato, Y. Katayama, T. Nishiya, T. Maeda, Technical Report of IEICE, 2002-03, pp. 57-64. Therein, an optical storage system with 2D modulation on a square lattice, with $d=1$ RLL constraints both in horizontal and vertical directions is described. For this system, a receiver, consisting of a 2D-equalizer and a 2D-Viterbi detector or a 2D-PRML detector, is disclosed in "Two-Dimensional Partial-Response Equalization and Detection Method with Multi-Track", T. Kato, S. Taira, Y. Katayama, T. Nishiya, T. Maeda, Technical Report of IEICE, 2002-03, pp. 65-70. The 2D PRML detector is based on three successive bit rows, but the typical add-compare-select operation (ACS) of the Viterbi-algorithm uses the HF-samples of the central bit row only; the other two bit rows are used in order to determine, in a joint way, the reference level from which the received HF-signal should be subtracted in order to derive the branch metrics for the branches

(or transitions) in the trellis diagram of the Viterbi detector. In this way, at its output, the Viterbi-detector yields bit-decisions for the central bit row only. In this sense, for successive rows, the PRML detectors operate already independently, and the state-
5 complexity for the complete set of bit rows has been reduced down to the complexity that is to be associated with 3 rows only. Within a strip of 3 rows, the known bit detector performs a kind of 2D-PRML, but with a 1D-output (for the single row being the mid-row of the 3-row strip). It should be noted that the channel strips are
10 processed independently, but that the state-complexity of the Viterbi-detector is still quite high.

[0007] Assuming the practical case of a 3-taps response in the tangential direction, as disclosed in the above mentioned documents. For the square lattice, but also when applying this
15 algorithm for a hexagonal lattice, assuming no modulation coding for both lattices, states characterized by 6 bits each would be obtained, yielding a number of $2^6=64$ states; each state would then have $2^3=8$ possible predecessors. On the square lattice, assuming the runlength modulation coding of the above references with 2D $d=1$
20 constraint, the number of states is only a bit smaller than 64, since some of the states are forbidden just because of the use of the runlength constraints along vertical and horizontal directions.

SUMMARY OF THE INVENTION

[0008] It is an object of the present invention to provide a Viterbi bit detection method that provides a high recording density, in particular, such that the traditional "eye" of the eye pattern may even be closed. The "eye height" in the traditional eye-pattern corresponds with the systematic minimum difference in signal levels for the case that a bit has a value "0" and the case that a bit has a value "1". An "open eye" means that (on average, or without any noise) the signal levels for bit "0" and bit "1" can be clearly discriminated: in such case, a threshold detection procedure with an appropriately set slicer level could be used. The case of a "closed eye" corresponds to the situation where some of the signal levels cannot unambiguously be allocated to bit "0" or bit "1", even in the absence of noise. There is, in the latter case, a range of signal levels, called the erasure zone, where the signal levels for bit "0" and bit "1" overlap.

[0009] It is a further object of the present invention to achieve a low bit error rate, which is particularly less than 10^{-2} to 10^{-1} as would be achieved for the case of a "closed eye" by application of a straightforward threshold detection prior to ECC decoding. Preferably, the symbol or byte error rate (BER) for "random" errors (as opposite to so-called "burst errors") in the case of a byte-oriented ECC, like the picket-ECC as used in BD (formerly known as DVR), shall not be larger than 2×10^{-3} ; for an uncoded channel bit stream this corresponds to an upper bound on the allowable channel-bit error rate (bER) of 2.5×10^{-4} .

[0010] Still further, a further reduction of the state-complexity of the independent Viterbi-detectors shall be achieved.

[0011] These objects are achieved according to the present invention by a Viterbi bit detection method ~~as claimed in claim 1,~~
5 according to which the channel data stream is stored on a record carrier along an N-dimensional channel tube, N being at least two, of at least two bit rows one-dimensionally evolving along a first direction and being aligned with each other along at least a second of N-1 other directions, said first direction together with said N-
10 1 other directions constituting an N-dimensional lattice of bit positions, comprising application of a row-based one-dimensional Viterbi bit detection method independent for each of the bit rows of said channel tube, wherein:

~~-~~[0012] calculation of the branch metrics for all possible state
15 transitions in a Viterbi trellis of a one-dimensional row-based Viterbi detector, said transitions representing a number of subsequent bits in said bit row, said bits being the central-row bits of a cluster of the N-dimensional lattice of bits, is based on the difference of the received HF signal value with respect to a
20 reference level, wherein said reference level depends on all bits of said cluster, said cluster comprising in addition to the central-row bits a number of primary ~~neighbouring~~ neighboring bits in each of a number of ~~neighbouring~~ neighboring bit rows on each side along said N-1 other directions of said central bit row along
25 which the one-dimensional Viterbi bit detection method is applied,

and wherein preliminary bit decisions for the primary ~~neighbouring~~
~~neighboring~~ bits in the ~~neighbouring~~~~neighboring~~ bit rows are used
for determining the reference level to be used for calculating said
branch metrics, and

5 -[0013] selection of the bit value for the central bit of said
cluster of the N-dimensional lattice of bits, corresponding with
said received HF signal value, is based on the calculated branch
metrics.

[0014] These objects are further achieved by a row-based Viterbi
10 bit detector ~~as claimed in claim 22~~ comprising Viterbi bit
detection units including means for calculation of the branch
metrics and means for selection of a bit value. The invention
relates further to a method of reproduction of a user data stream,
which is error correction code encoded and modulation code encoded
15 into a channel data stream and stored on a record carrier,
comprising a Viterbi bit detection method as described above for
detecting the bit values of bits of the channel data stream and a
modulation code decoding method and an error correction code
decoding method. Still further, the present invention relates to a
20 reproduction device ~~as claimed in claim 25~~ and a computer program
~~as claimed in claim 26. Preferred embodiments of the invention are~~
~~defined in the dependent claims.~~

[0015] The present invention is based on the idea to achieve
reliable bit detection by using a number of independent 1D Viterbi
25 sequence-detectors, one for each bit row in the channel tube: the

interference between successive ~~neighbouring-neighboring~~ bit rows is taken into account via the computation of the branch metrics for the considered bit row, in which preliminary bit decisions on the primary ~~neighbouring-neighboring~~ bits are used that may require
 5 local bit decisions on the secondary ~~neighbouring-neighboring~~ bits in the ~~neighbouring-neighboring~~ bit rows, said secondary ~~neighbouring-neighboring~~ bits being the ~~neighbouring-neighboring~~ bits of said primary ~~neighbouring-neighboring~~ bits not being part of the central bit row being considered for the 1D Viterbi
 10 detection.

[0016] Regarding the general layout and function of a PRML bit detection apparatus, reference is made to the ~~above-above~~-mentioned WO 00/18016 where also several terms are explained. This description and explanation is herein incorporated by reference.

15 [0017] ~~Preferred-In preferred~~ embodiments for determining the preliminary bit decisions on the primary ~~neighbouring-neighboring~~ bits ~~are defined in claims 2 to 7. Thus,~~ a slicer level can be used in a threshold detection. Further, specific bit values of the bits in the central row can be used in accordance with each of the
 20 specific branches to be considered in the Viterbi trellis. Said threshold detection is based on the detected HF signal value for a particular bit without consideration of the HF-signal samples at the ~~neighbouring-neighboring~~ bit-locations.

[0018] According to a preferred embodiment, a predetermined
 25 criterion is evaluated for determining the preliminary bit

decisions, ~~which this criterion is being~~ determined by the sum over all the primary ~~neighbouring-neighboring~~ bits, said sum comprising terms related to a subcriterion that is based on the differences of the HF signal value and a reference HF signal value corresponding to the bit cluster of each single primary ~~neighbouring-neighboring~~ bit, ~~which the evaluation is being~~ done for all possible bit units obtained for all possible values of said primary ~~neighbouring-neighboring~~ bits, and wherein the bit unit with the lowest value of said predetermined evaluation criterion is selected. Preferred subcriteria ~~which~~ relate to the squared value or the absolute value of the difference of the HF signal value and a reference HF signal value corresponding to the bit cluster of each single primary ~~neighbouring-neighboring~~ bit ~~to are defined in claims 5 and 6.~~

[0019] In addition to the primary ~~neighbouring-neighboring~~ bits, preliminary bit decisions on secondary ~~neighbouring-neighboring~~ bits can be used for determining the preliminary bit decisions on the primary ~~neighbouring-neighboring~~ bits ~~as defined in claim 8.~~ Those preliminary bit decisions on said secondary ~~neighbouring-neighboring~~ bits can be obtained, for instance, by threshold detection.

[0020] There are different ways to calculate the branch metrics. ~~Two preferred ways are defined in claims 9 and 10.~~ The Viterbi algorithm in a PRML bit detector searches for the "best" path, that is, the path with the minimum path cost. The path cost is sometimes called "path metric". A path can be seen as a succession of

transitions. A transition from one state to another state is called a branch. Each transition (or branch) has associated with it a certain branch metric (or branch cost). The path metric for a given path is the sum of the costs of the individual branches of the path, that is, the path metric is a sum of a selection of branch metrics.

[0021] Generally, the present invention is applicable to a multi-dimensional code, where the channel words of the channel data stream may evolve in more than one direction as is the case for a card-based system, i.e., where the channel data stream is stored on a record carrier along a multi-dimensional channel tube with dimension at least two. Therein, the first direction along which the bit rows evolve is preferably common to all bit rows of the channel tube. The first direction constitutes, together with the N-1 other directions, along which the bit rows are aligned with each other, an N-dimensional space and an N-dimensional lattice of bit positions. The channel tube comprises at least two bit rows of channel bits evolving along said first direction, and the collection of all said channel tubes fill the whole N-dimensional space.

[0022] However, it is preferred to apply the invention to a channel data stream which comprises a one-dimensionally evolving bit sequence, or which comprises a channel strip of at least two bit rows one-dimensionally evolving along a first direction and aligned with each other along a second direction, preferably

oblique or even orthogonal to said first direction, said two directions constituting a two-dimensional lattice of bit positions. Preferred embodiments of such a lattice are a 2D lattice of a square or a hexagonal type ~~as defined in claims 12 and 13.~~

5 [0023] In a hexagonal lattice, hexagonal clusters may be formed of a set of 7 bits in total, comprising three bits in a central bit row and two primary ~~neighbouring~~ neighboring bits in each of the two ~~neighbouring~~ neighboring bit rows. Further secondary ~~neighbouring~~ neighboring bits can be located in the ~~neighbouring~~ neighboring bit rows of the central bit row considered. Preferred
10 embodiments of the invention ~~using~~ use hexagonal clusters ~~are defined in claims 14 to 18.~~

[0024] An advantageous embodiment for calculation of the branch metrics ~~using~~ uses an expectation value ~~is defined in claim 19.~~ The
15 method can also be applied in the three-dimensional case where the bits are location on bit positions of a three-dimensional lattice ~~as claimed in claim 20.~~

[0025] The bit detection method according to the invention can also include an iterative use of the row-based one-dimensional
20 Viterbi bit detection method: the output of the 1D-Viterbi detectors for a given 1D part of the set of bit rows can be used for the required primary bit decisions in the ~~neighbouring~~ neighboring rows during a second run of the method for the same 1D part (same bits along the 1D row) of the set of rows. The purpose
25 is to use the output of the first set of 1D-Viterbi detectors for

all bit rows as a better bit decision for the primary bit decisions required in a possible second set of Viterbi detectors for all bit rows.

5

BRIEF DESCRIPTION OF THE DRAWINGS

[0026] The invention shall now be explained with reference to the drawings, in which:

[0027] Fig. 1 shows a block diagram of a general layout of a coding system_{7i}

10 [0028] Fig. 2 shows a schematic diagram indicating a strip-based two-dimensional coding scheme_{7i}

[0029] Fig. 3 shows a schematic signal-pattern for a two-dimensional code on hexagonal lattices_{7i}

15 [0030] Fig. 4 shows a row-based partitioning of the 2D target response on the hexagonal bit cluster_{7i}

[0031] Fig. 5 shows a Finite-State-Machine for an 1D-PRML with 3-tap PRML target_{7i}

[0032] Fig. 6 shows a hexagonal bit-cluster with enumeration convention used in this application_{7i}

20 [0033] Fig. 7 illustrates the invention in case two primary ~~neighbouring~~ neighboring bits are used_{7i}

[0034] Fig. 8 illustrates the invention in case three primary ~~neighbouring~~ neighboring bits are used_{7i}

[0035] Fig. 9 shows a trellis for 1D-PRML bit detection_{7i}

[0036] Fig. 10 shows a repeated trellis for 1D-PRML bit detection_i;

[0037] Fig. 11 illustrates a particular example using the trellis shown in Fig. 9_i;

5 [0038] Fig. 12 illustrates the determination of the path cost according to the present invention_i;

[0039] Fig. 13 shows a block diagram of a bit detector in case two primary ~~neighbouring~~_i ~~neighboring~~ bits are used in each primary ~~neighbouring~~_i ~~neighboring~~ bit row_i;

10 [0040] Fig. 14 shows the HF reference signal levels used in the bit detector shown in Fig. 13_i;

[0041] Fig. 15 shows a single bit detection unit of a bit detector shown in Fig. 13_i;

15 [0042] Fig. 16 shows a Fermi-Dirac like S-curve for deriving soft-decision information from a HF-signal of a bit_i;

[0043] Fig. 17 shows reference levels derived from the signal pattern_i and

[0044] Fig. 18 shows the bit-error-rate as a function of SNR (defined relative to the full-reflection signal level) for a
20 typical density.

DESCRIPTION OF THE PREFERRED EMBODIMENTS

[0045] Fig. 1 shows typical coding and signal processing elements of a data storage system. The cycle of user data from
25 input DI to output DO can include interleaving 10, error-

correction-code (ECC) and modulation encoding 20, 30, signal preprocessing 40, data storage on the recording medium 50, signal post-processing 60, binary detection 70, and decoding 80, 90 of the modulation code, and of the interleaved ECC. The ECC encoder 20
5 adds redundancy to the data in order to provide protection against errors from various noise sources. The ECC-encoded data are then passed on to a modulation encoder 30 which adapts the data to the channel, i.e., it manipulates the data into a form less likely to be corrupted by channel errors and more easily detected at the
10 channel output. The modulated data are then input to a recording device, e.g., a spatial light modulator or the like, and stored in the recording medium 50. On the retrieving side, the reading device (e.g., photo-detector device or charge-coupled device (CCD)) returns pseudo-analog data values which must be transformed back
15 into digital data (one bit per pixel for binary modulation schemes). The first step in this process is a post-processing step 60, called equalization, which attempts to undo distortions created in the recording process, still in the pseudo-analog domain. Then, the array of pseudo-analog values is converted to an array of
20 binary digital data via a bit detector 70. The array of digital data is then passed first to the modulation decoder 80, which performs the inverse operation to modulation encoding, and then to an ECC decoder 90.

[0046] In the ~~above-mentioned~~ European ~~patent~~-Patent
25 ~~application~~-Application No. EP 01203878.2, the 2D constrained

coding on hexagonal lattices in terms of nearest-~~neighbour~~-neighbor clusters of channel bits is described. Therein, it has been ~~focussed~~-focused mainly on the constraints with their advantages in terms of more robust transmission over the channel, but not on the actual construction of such 2D codes. The latter topic is addressed in ~~the European patent~~-Patent application-Application No. 02076665.5 (PHNL 020368), i.e., the implementation and construction of such a 2D code is described therein. By way of example, a certain 2D hexagonal code shall be illustrated in the following.

However, it should be noted that the general idea of the invention and all measures can be applied generally to any 2D code, in particular, any 2D hexagonal or square lattice code. Finally, the general idea can also be applied to multi-dimensional codes, possibly with isotropic constraints, characterized by a one-dimensional evolution of the code.

[0047] As mentioned, in the following, a 2D hexagonal code shall be considered. The bits on the 2D hexagonal lattice can be identified in terms of bit clusters. A hexagonal cluster consists of a bit at a central lattice site, surrounded by six nearest ~~neighbours~~-neighbors at the ~~neighbouring~~-neighboring lattice sites. The code evolves along a one-dimensional direction. A 2D strip consists of a number of 1D rows, stacked upon each other in a second direction orthogonal to the first direction. The principle of strip-based 2D coding is shown in Fig. 2. Between the strips or

between groups of successive strips, a guard band of, for instance, one row may be located.

[0048] The signal-levels for 2D recording on hexagonal lattices are identified by a plot of amplitude values for the complete set of all hexagonal clusters possible. Use is further made of the isotropic assumption, that is, the channel impulse response is assumed to be circularly symmetric. This implies that, in order to characterize a 7-bit cluster, it only matters to identify the central bit, and the number of "1"-bits (or "0"-bits) among the nearest-~~neighbour~~-neighbor bits (0, 1, ..., 6 out of the 6 ~~neighbours~~-neighbors can be a "1"-bit). A "0"-bit is a land-bit in our notation. A typical "Signal-Pattern" is shown in Fig. 3. Assuming a broad-spiral consisting of 11 parallel bit rows, with a guard band of 1 (empty) bit row between successive broad spirals, the situation of Fig. 3 corresponds to a density increase with a factor of 1.7 compared to traditional 1D optical recording (as used in e.g., in the Blu-ray Disc (BD) format (using a blue laser diode)).

[0049] According to the present invention, the broad spiral (or meta-spiral) consists of a number of bit rows. It is proposed to apply a row-based 1D-PRML, in which the Viterbi-trellis relates only to the bits in the individual bit row itself. For a 3-tap target response in the direction along the bit row, states are obtained that are each defined by two bits. Fig. 4 shows the 2D target response on the cluster of bits of the hexagonal lattice.

The corresponding Finite-State Machine (FSM) is shown in Fig. 5.

The FSM reveals the syntax when going from one state "i" (with bits $(b_0^i b_1^i)$) towards another state "j" (with bits $(b_0^j b_1^j)$). Such a transition between states "i" and "j" is only allowed on the

5 condition that $b_1^i = b_0^j$, i.e., the second bit of the first state and the first bit of the second state must be identical. It should be noted that a transition from one state "i" to another state "j" completely characterizes the three bits of the 7-bits cluster at the central bit row by $(b_0^i b_1^i b_1^j)$. No further restrictions are
10 present on the possible transitions, because in the practical description made here, there is no 1D RLL constraint assumed for the case of our 2D hexagonal-lattice modulation.

[0050] Next, the computation of branch metrics shall be explained. Fig. 6 shows the enumeration of the 7 bits of an
15 hexagonal cluster: $x_0, x_1, x_2, x_3, x_4, x_5$, and x_6 . The branch metric for the transition from state "i" to state "j" is denoted by β_{ij} . It is preferably given by:

$$\beta_{ij} = \left(HF_0 - R.L. [x_0 = b_1^i; x_1 = b_1^j; x_2; x_3; x_4 = b_0^j; x_5; x_6] \right)^2$$

20

where HF_0 denotes the sample value of the (possibly equalized) received signal at the (central) bit x_0 . R.L. denotes the reference amplitude level, which depends on all the bit-values of the 7-bit cluster. For a given transition ("i" to "j"), the values of the
25 bits x_0, x_1 and x_4 in the central bit row are already fixed. The

other bits that still need to be determined, occur in two pairs of bits, denoted by x_2 , x_3 and x_5 , x_6 . These pairs correspond with nearest ~~neighbour~~-neighbor bits of the central bit x_0 , where each of these pairs is located either in the upper ~~neighbouring~~ neighboring bit row, or in the lower ~~neighbouring~~-neighboring bit row. The bit values in these two bit-pairs are required in order to be able to uniquely identify the reference level R.L. to be used in the branch metric for the transitions in the 1D-PRML detector for the considered bit row. These bits will be further referred to as the primary ~~neighbouring~~-neighboring bits. Therefore, these bit decisions on the primary ~~neighbouring~~-neighboring bits can be seen as preliminary bit decisions, needed to assist the evaluation of the branch metrics required for the bit decisions of the bit in the considered (central) bit row. One aspect of the present invention relates to the decisions that are to be made on those two bit-pairs of primary ~~neighbouring~~-neighboring bits. The quality of these (temporarily needed) preliminary bit decisions that are (only) required for the computation of branch metrics, have an impact on the quality of the ultimate bit decisions of the bits in the central row on which the 1D-PRML is applied. The preliminary bit decisions on the primary ~~neighbouring~~-neighboring bits are thus never used as real output bits for the ~~neighbouring~~-neighboring bit rows.

[0051] A rather straightforward approach is to use threshold detection for the primary ~~neighbouring~~-neighboring bits denoted by

x_2 and x_3 and x_5 and x_6 .- However, threshold detection is highly unreliable due to the large overlap in signal levels in the signal pattern shown in Fig. 3. In the performance analysis, it turns out that not much performance gain is obtained from the use of 1D-PRML with threshold detection for the primary ~~neighbouring-neighboring~~ bits compared to threshold detection applied as the bit detector on the whole meta-spiral. The quality of the primary ~~neighbouring-neighboring~~ bit decisions is obviously not good enough.

[0052] More reliable bit decisions at the primary ~~neighbouring-neighboring~~ bits can be obtained through the use of the hard-decision bit detectors on nearest-~~neighbour-neighbors~~ bits in ~~neighbouring-neighboring~~ bit rows.

[0053] Fig. 7 illustrates the use of a HD-2 (HD = hard-decision) bit detector which uses preliminary bit decisions on a doublet of two primary ~~neighbouring-neighboring~~ bits in each of the upper and lower ~~neighbouring-neighboring~~ bit row of the central bit row. As explained above, the bit values of the bits in the central bit row are fixed for a given branch in the trellis. The bit-values of the primary ~~neighbouring-neighboring~~ bits in the ~~neighbouring-neighboring~~ bit rows are determined as follows, wherein the bit-pairs (also referred to as bit units) denoted by x_2 and x_3 , and x_5 and x_6 are treated independent of each other: a criterion, preferably the sum of two terms, each term being a square of a difference of the actual received HF sample at one bit of the bit pair with the corresponding reference level for that primary ~~neighbouring~~

neighboring bit in the bit pair, is evaluated for all possible two-bit combinations of the bits in the bit pair. The bit-pair which has the lowest value for the criterion is selected. This yields the preliminary bit values to be used in the reference level for the branch metric at the bit x_0 in the central row. But also for the primary ~~neighbouring~~-neighboring bit values, reference levels need to be specified for the criterion of the HD-2 detector. Therefore, the bit values of all nearest-~~neighbour~~-neighbor bits for each bit of the two bit pairs of primary ~~neighbouring~~-neighboring bits are needed as well. Some of these ~~neighbours~~-neighbors are fixed by the branch considered, others are part of the HD-2 selection procedure; all other remaining nearest-~~neighbour~~-neighbor bits for each primary ~~neighbouring~~-neighboring bit of the two bit-pairs are referred to as secondary ~~neighbouring~~-neighboring bits, which may preferably be obtained by threshold detection (TD).

[0054] Fig. 8 illustrates the use of a HD-3 bit detector which uses preliminary bit decisions on triplets of three primary ~~neighbouring~~-neighboring bits in the two upper and the two lower ~~neighbouring~~-neighboring bit rows of the central bit row. The procedure is quite similar to the one explained above for the HD-2 bit detector. Preliminary bit decisions are required at the primary ~~neighbouring~~-neighboring bits of the two bit pairs denoted by x_2 and x_3 , and x_5 and x_6 . Each bit pair is again processed independently. To each bit pair is added a third bit in the second ~~neighbouring~~-neighboring bit row of the central bit row so that a

bit-triplet (also referred to as bit unit) is formed. There are 8 possible bit-triplets. As above, a criterion is evaluated for each of these 8 possibilities, and the one with the lowest value is selected. The criterion used for the HD-3 bit detector is a sum of
5 three terms, one for each bit in the triplet. Each term corresponds to one bit in the triplet, and preferably is the squared value of the difference between the measured HF signal at that bit and the corresponding reference level. The latter is determined by using bit values of the nearest-~~neighbour~~-neighbor bits of the bits in
10 the triplets; some of these nearest ~~neighbour~~-neighbor bits are fixed by the branch considered, others are part of the HD-3 selection procedure, and yet other bits, i.e., the 6 bits surrounding the bit-triplet in the ~~neighbouring~~-neighboring bit rows of the central bit row, which bits are further referred to as
15 secondary ~~neighbouring~~-neighboring bits, may preferably be determined by threshold detection (TD). From the chosen bit-triplet, only the bits of the bit pairs x_2 and x_3 for the bottom bit-triplet, and x_5 and x_6 for the top bit-triplet are needed in order to choose the reference level that has to be used in the
20 computation of the branch metric for the 1D-PRML at the bit x_0 in the considered bit row (assuming that the intersymbol-interference of which the bit detector can take care ~~off~~-of resides within a hexagonal cluster comprising not more than 7 bits). However, inclusion of the extra third bit in each of both 3-bit bit units
25 may largely improve the quality of the preliminary bit decisions on

the primary ~~neighbouring~~-neighboring bit pairs x_2 and x_3 for the bottom bit-triplet, and x_5 and x_6 for the top bit-triplet.

[0055] The trellis for 1D-PRML bit detection with 3-taps impulse-response in the tangential direction is illustrated in Fig.

5 9. As can be seen, each state has exactly two predecessor states.

[0056] These two predecessor states are the states that have, as a last bit, the first bit of the current (considered) state. For instance, transitions from state "01" are only allowed for the states "10" and "11" as next state. These transitions yield the 3-bit sequences "010" and "011", respectively.

[0057] Fig. 10 shows the repeated trellis for 1D-PRML bit detection with 3-taps impulse-response in the tangential direction.

[0058] All paths through the trellis realize all possible bit sequences. The Viterbi algorithm (for maximum-likelihood sequence detection) searches for the "best" path, that is, the path with the minimum path cost. The path cost is sometimes called "path metric". A path can be seen as a succession of transitions. A transition from one state to another state (from moment k towards moment $k+1$) is also called a branch. Each transition (or branch) has associated with it a certain branch metric (or branch cost). The path metric for a given path is the sum of the costs of the individual branches of the path, that is, the path metric is a sum of branch metrics.

[0059] For the above case (with 2-bits states "00", "01", "10" and "11"), a branch metric for a transition between states s_0 and s_1 , from moment (or time) $k-1$ to moment k , is the squared value for

the L2-norm (or absolute value for the L1-norm) of the difference between the measured HF-sample at time k (denoted by HFk) and the reference level R.L. that is associated with the transition from state s0 to state s1. The reference level is a kind of ideal

5 (noise-free) signal level for the transition that is considered.

[0060] The reference level here also depends on the bit values of the ~~neighbouring-neighboring~~ bit rows, because of the strong 2D inter-symbol-interference (ISI). In standard 1D-storage, the ~~neighbouring-neighboring~~ bit rows are always far away so that this
10 problem is not present there.

[0061] The applied Viterbi algorithm shall now be explained with reference to Fig. 11. Finding the best path via a brute force (exhaustive) search is not favorable because of the computational burden. Viterbi has introduced the procedure of "dynamic

15 programming" with a complexity that is linearly growing with the length of the bit-sequence to be determined. In the following, it shall be assumed that the best path at time k that arrives in state "01" shall be found. It is further assumed that the following two aspects for each of the states have previously been evaluated at
20 time k-1:

(a) for each state, the path metric for the best path, i.e., the path with minimum cost, that arrives in that state, is known;

(b) for each state, the predecessor state at the previous time is known: the predecessor state is the state that lies on the
25 best path at the previous time moment.

[0062] Then, the best path at moment k that arrives in "01" can be found by looking at the two possible predecessor states of state "01": one path starts in state "00" with a path metric $p_{00,k-1}$ and needs a branch "00"→"01" to arrive in state "01", with a branch
 5 metric given by $b_{00 \rightarrow 01,k}$. The second path starts in state "10" with a path metric $p_{10,k-1}$ and needs a branch "10"→"01" to arrive in state "01", with a branch metric given by $b_{10 \rightarrow 01,k}$. The best path is the path that realizes the minimum cost:

10 minimum of $p_{00,k-1} + b_{00 \rightarrow 01,k}$ and $p_{10,k-1} + b_{10 \rightarrow 01,k}$.

If the latter is the best path, then the predecessor state of state "01" is "10", in the other case, the predecessor state of state "01" is "00". This procedure is known as Add-Compare-Select (ACS):
 15 the branch metric is added to the previous path metric to obtain a candidate for the current path metric; the two candidate path metrics are compared, and the path with the lowest metric is selected.

[0063] This procedure is repeated for all states (at each time
 20 moment) .. As a result, a collection of path metric and predecessor state for each state and each time moment is obtained. The actual bit detection is then performed by a so-called back-tracking operation: one starts from a best state at moment k , and goes back to its predecessor state at time $k-1$, and to the predecessor state
 25 of that state at moment $k-2$, and so on. The back-tracking is done

for a certain depth K (known as back-tracking depth); the bit value at moment $k-K$ is (for instance) the bit-value of the first bit in the resulting state obtained at the end of the back-tracking operation. States "00" and "01" lead to bit "0"; states "10" and
 5 "11" lead to bit "1".

[0064] There are four states, denoted by "00", "01", "10" and "11". For each of these states, the best predecessor of that state and the path costs up to the given state for the path with the lowest cost going to that state are evaluated, as illustrated in
 10 Fig. 12. For state "00", for instance, possible predecessor states are the states "00" and "10", leading to the bits "000" (from state "00" towards state "00") and "100" (from state "10" towards state "00") for the triplet of bits denoted by (b_4, b_0, b_1) , as denoted by the upper and lower line related to state "00". These separate
 15 transitions towards state "00" are called branches of the trellis. The bit b_0 is the central bit of the triplet. For a given state, thus two possible branches are there. For each of the branches, the most likely candidates for the bits b_2 and b_3 in the lower row relative to the considered row, and for the bits b_5 and b_6 in the
 20 upper row relative to the considered row are determined.

[0065] According to the embodiment shown in Fig. 12, this is done by using an HD-2 detector for each of the two bit-pairs $(b_2$ and $b_3)$ and $(b_5$ and $b_6)$: the respective HD-2 detectors are denoted by "HD-2 Upp" and "HD-2 Low" in Fig. 12. The resulting four bits
 25 $(b_2$ and b_3, b_5 and $b_6)$ together with the bits of the bit-triplet

(b_4, b_0, b_1) define the 7 bits of the hexagonal cluster: these 7 bits uniquely define an index for the specific reference level in the memory of Reference Levels, the latter being denoted by HF Ref.Lev.Mem. The reference levels for each of the two branches
 5 (for each of the four states) is compared with the actually measured HF signal for bit b_0 , denoted by HF. Such comparison can be done by a squared value in case of the L2-norm (or absolute value in case of the L1-norm) of the difference of the received signal and the reference level: such difference value yields the
 10 actual branch metric for each of the two transitions in the trellis. A standard Add-Compare-Select (ACS) unit that uses the current branch metrics and the path metrics up to the two possible preceding states, i.e., "00" and "10" in the case of state "00", further determines the best predecessor of the current state,
 15 denoted as pre_{00} , and the path cost for the cheapest path up to the currently considered state "00", denoted by $paco_{00}$. This procedure is done for each of the four possible states separately and independently, since no information exchange is needed between the distinct procedures for each of the four states.
 20 **[0066]** Fig. 13 shows the HD-2 bit detector HD-2 Upp for the determination of the upper-row bits (b_5 and b_6). The HD-2 Upp block has as input 8 bit values of the ~~neighbouring~~neighboring bits, three of which are set by the considered branch in the trellis for the central row, which is the row under consideration; the other 5
 25 bits are referred to as secondary ~~neighbouring~~neighboring bits,

and are obtained as simply ~~threshold~~-threshold-detected bits derived from the HF-samples at the corresponding bit-positions. The HD-2 Upp block also has two HF-samples as input for the bit-locations of bits b_5 and b_6 .

5 [0067] The so-called "branch-bits" apply for the three bits of a given transition in the trellis for the bit row under consideration.

[0068] A similar diagram (not shown) applies for the block HD-2 Low, yielding bit decisions in the lower row for the bits denoted
10 by b_2 and b_3 . The bits determined by the HD-2 Upp block are further used, together with the similarly derived - by means of block HD-2 Low - lower bits b_2 and b_3 to derive reference levels from a reference level memory as shown in Fig. 12.

[0069] Reference levels to be used in the core of HD-2 bit
15 detector are shown in Fig. 14. The bit-numbering applied in Figs. 14 and 15 should be noted which refers to the order of the bits in the bit pair. For the first bit in the bit pair, bit b_0 , the reference levels are denoted with the first sub-script underlined; for the second bit in the bit pair, bit b_1 , the reference levels
20 are denoted with the second sub-script underlined.

[0070] Fig. 15 shows the basic layout of the HD-2 bit detector block. This block describes both the HD-2 Upp and the HD-2 Low blocks. Inputted are 8 nearest ~~neighbouring~~-neighboring bits on the hexagonal lattice and HF-samples of the two bits of the bit pair
25 that needs to be updated. Outputted are the two updated bit values,

i.e., the HD-2 detected bits. For each hexagonal cluster with one central bit and 6 ~~neighbouring-neighboring~~ bits, a reference signal level is available from a memory, i.e., the HF Reference-Level Memory. The reference level to be taken from the memory is

5 determined by the two bits of the bit pair, and by 5 out of the 8 ~~neighbouring-neighboring~~ bits of the bit pair. The 8 ~~neighbouring-neighboring~~ bits of the bit pair comprise 3 bits of the central bit row (determined by the actual branch considered), and 5 secondary ~~neighbouring-neighboring~~ bits.

10 [0071] The received HF-signal for each bit of the two bits of the bit pair are subtracted from the corresponding reference level; the absolute values (shown here; it may also be any other "norm" like the quadratic norm using the squared values instead) of these respective signal differences are added together for each of the

15 four possible two-bit configurations for the two bits of the bit pair. The bits that result from the HD-2 bit detector are those that lead to the smallest value of the above set of 4 parameters or samples of the selection criterion, one sample for each possible bit pair. This is denoted in Fig. 15 in short-hand notation by arg

20 min: the arguments (bits b_0 and b_1 of the bit pair) for which the criterion is at the minimum.

[0072] According to an alternative embodiment, soft-decision information about the primary ~~neighbouring-neighboring~~ bits in the bit pairs (denoted by x_2 and x_3 , and x_5 and x_6) is used. The branch

25 metric for a given transition from state "i" to state "j" is then

computed as an expectation value, which is the average taken over all possible bit-configurations in the two bit pairs of primary ~~neighbouring~~ neighboring bits. Formally, this can be written as (with the index k to the HF_k signal referring to the sample at the k-th bit in the hexagonal cluster):

$$\beta_{ij} = \sum_{b_2=0}^1 \sum_{b_3=0}^1 \sum_{b_5=0}^1 \sum_{b_6=0}^1 p(x_2=b_2; x_3=b_3; x_5=b_5; x_6=b_6 \mid x_0=b_1^i; x_1=b_1^j; x_4=b_0^i; HF_2, HF_3, HF_5, HF_6) \\ (HF_0 - R.L.[x_0=b_1^i; x_1=b_1^j; x_2=b_2; x_3=b_3; x_4=b_0^i; x_5=b_5; x_6=b_6])^2$$

10

[0073] It should be noted that the bit values denoted by (b_0^i, b_1^i) refer to the two bits in state "i" of the 4-state Viterbi-trellis (shown in Fig. 9), and similarly for state "j". The probability factor in the above expression can be split into separate factors for each of the two independent bit pairs. For each bit pair, the factor can further be split into factors relating to the individual bits, yielding:

15

$$p(x_2=b_2; x_3=b_3; x_5=b_5; x_6=b_6 \mid x_0=b_1^i; x_1=b_1^j; x_4=b_0^i; HF_2, HF_3, HF_5, HF_6) = \\ p(x_2=b_2 \mid \text{all 6 nearest neighbours of } x_2; HF_2) \times \dots \times \\ p(x_6=b_6 \mid \text{all 6 nearest neighbours of } x_6; HF_6)$$

20

Just as in the HD-2 bit detector, reference can be made to Fig. 7, but now primary ~~neighbouring~~ neighboring bits of the bit pairs refer to bits with soft-decision information. The nearest-~~neighbour~~ neighbor bits of the bit pairs - those not part of the central bit

25

row, nor from the HD-2 bit-unit are referred to as secondary
~~neighbouring-neighboring~~ bits - are determined by threshold
 detection. For each bit of the bit pairs, all of its nearest-
~~neighbour-neighbor~~ bits are thus characterized. Soft-decision
 5 information can, for instance, be determined by a "Fermi-Dirac"-
 like S-curve as shown in Fig. 16 (given the configuration
 ("config") of nearest-~~neighbour-neighbor~~ bits, and the HF sample).
 Therein, Fermi-Dirac-like S-curves for deriving soft-decision
 information from the HF-signal of a bit at position (k, l) are
 10 shown. T_0 is the reference level when the central bit (at (k, l))
 is a zero, T_1 applies for the case that the central bit is a one.
 The different curves relate to different noise variances. The
 reference levels T_0 and T_1 are derived from the signal pattern as
 shown in Fig. 17 where an example is given when two nearest
 15 ~~neighbour-neighbor~~ bits in the hexagonal cluster are equal to "1".

[0074] The performance for various detectors for the density of
 1.4x the density of BD has been computed. A lattice parameter $a =$
 165nm with a pit-hole diameter equal to 120nm (in order to avoid
 signal folding) has been assumed. The channel is subjected to AWGN
 20 disturbance (additive white Gaussian noise). The detectors are:

- threshold detection (TD);
- HD-3 Hard-Decision Iterative Bit detector (HD-3);
- TD-assisted 1D-PRML;
- HD-2 assisted 1D-PRML;
- 25 - HD-3-assisted 1D-PRML;

- SD-1 Soft-Decision Iterative Bit detector (SD-1);
- Soft-Decision-assisted 1D-PRML.

The results are shown in Fig. 18. Evaluation is done in terms of the (channel) bit-error-rate (bER), as a function of the SNR of the channel. It should be noted that the TD-assisted 1D-PRML has a high bER, that is, that version of the 1D-PRML detector offers only a marginal improvement in bER compared to the TD detector itself. The HD-3-assisted 1D-PRML, on the other hand, is almost identical in performance to the soft-decision bit detector SD-1, and slightly better than the hard-decision HD-3 bit detector. The HD-3 assisted 1D-PRML is even better than the SD-assisted 1D-PRML: this might be due to the fact that the soft-decision information is obtained per bit only (and uses TD-decisions at some of its ~~neighbours~~neighbors), whereas, the HD-3 assisted 1D-PRML detector searches for optimum joint (hard) bit-decisions in an area of 3-bits at each side of the considered track.

[0075] The present invention provides a solution to achieve reliable bit detection by using a number of independent 1D-Viterbi bit detectors (also known as sequence detectors) is used, one for each bit row in the channel tube: the interference between successive ~~neighbouring~~neighboring bit rows is taken into account via the computation of the branch metrics (for the considered bit row), in which local bit decisions on the primary ~~neighbouring~~neighboring bits in the ~~neighbouring~~neighboring rows are used. As local bit detectors going beyond the performance of a threshold

detector, the use of a HD-2 or HD-3-like hard-decision bit detector is proposed. Other local bit detectors might also be used, insofar that they take account of the specific bit values for the respective branches in the Viterbi trellis of the central row that is being processed with one-dimensional row-based Viterbi bit detector to condition the preliminary bit detection for the primary ~~neighbouring-neighboring~~ bits in the ~~neighbouring-neighboring~~ rows of the considered bit row.

[0076] Further, it is proposed to use the output of a soft-

decision bit detector at the bits in the ~~neighbouring-neighboring~~ rows of the central row, in order to compute the branch metrics.

Practically, it is proposed to use soft-decision information that can be directly generated from the signal pattern (with 2x7 signal levels, grouped as 7 pairs of levels). Of course, other soft-

decision bit detectors can be used for the same purpose, like e.g., iterative soft-decision-detectors. Preferably, the two- and three-dimensional cases are advantageous where the bits are arranged on a two- or three-dimensional lattice.

ABSTRACT+ OF THE DISCLOSURE

~~The present invention relates to a~~ Viterbi bit detection method for detecting the bit values of bits of a channel data stream stored on a record carrier along an N-dimensional channel tube, N being at least two, of at least two bit rows one-dimensionally evolving along a first direction and being aligned with each other along at least a second of N-1 other directions, ~~said the first direction together with said the N-1 other~~ directions constituting an N-dimensional lattice of bit positions, ~~comprising includes~~ application of a row-based one-dimensional Viterbi bit detection method independent for each of the bit rows of said channel tube. To achieve a reliable bit detection, a number of independent one-dimensional row-based Viterbi bit detectors, also known as sequence detectors, is used, one for each bit row in the channel tube: the interference between successive ~~neighbouring~~ neighboring bit rows is taken into account via the computation of the branch metrics (for the considered bit row), in which local bit decisions on the primary ~~neighbouring~~ neighboring bits in the ~~neighbouring~~ neighboring rows are used. As local bit detectors going beyond the performance of a threshold detector, the use of a HD-2 or HD-3-like hard-decision bit detector is proposed. ~~Other local bit detectors might also be used, insofar that they take account of the specific bit values for the respective branches in the Viterbi trellis of the central row that is being processed with~~

PHNL 020930

~~one-dimensional row-based Viterbi bit detector to condition the preliminary bit detection for the primary neighbouring bits in the neighbouring rows of the considered bit row.~~

5 ~~Fig. 7~~